A propositional proof system $R$ is called *automatizable* if there is an algorithm that produces for every formula $F$ a proof in $R$ in time polynomial in $s(F)$, where $s(F)$ is the length of the shortest proof of $F$ in $R$.

This paper shows that under a plausible complexity assumption, the proof systems *Polynomial Calculus* and *Polynomial Calculus with Resolution* are not automatizable. The same result was shown for the well-known *Resolution* proof system under the same complexity assumption by Alekhnovich and Razborov [1]. The proof closely follows the method of Alekhnovich and Razborov [1].

References